Higher-Order Label Homogeneity and Spreading in Graphs

Dhivya Eswaran
deswaran@cs.cmu.edu
Carnegie Mellon University

Srijan Kumar
srijan@gatech.edu
Georgia Institute of Technology

Christos Faloutsos
christos@cs.cmu.edu
Carnegie Mellon University

ABSTRACT
Do higher-order network structures aid graph semi-supervised learning? Given a graph and a few labeled vertices, labeling the remaining vertices is a high-impact problem with applications in several tasks, such as recommender systems, fraud detection, and protein identification. However, traditional methods rely on edges for spreading labels, which is limited by the fact that all edges are not equal. Vertices with stronger connections participate in higher-order structures in graphs, which calls for methods that can leverage these structures in the semi-supervised learning tasks.

To this end, we propose Higher-Order Label Spreading (HOLS) to spread labels using higher-order structures. HOLS has strong theoretical guarantees and reduces to standard label spreading [38] in the base case. Via extensive experiments, we show that higher-order label spreading using triangles in addition to edges (HOLS-3) is up to 4.7% better than label spreading using edges alone. Compared to prior traditional and state-of-the-art methods, the proposed HOLS-3 leads to statistically significant accuracy gains in all-but-one cases, while remaining fast and scalable to large graphs.

ACM Reference Format:

1 INTRODUCTION
Given an undirected unweighted graph and a few labeled vertices, the graph transductive learning or semi-supervised learning (SSL) aims to infer the labels for the remaining unlabeled vertices [1, 6, 17, 29, 35, 36, 38, 39]. Graph SSL finds applications in a number of settings: in a social network, we can infer a particular characteristic (e.g. political leaning) of a user based on the information of her friends to produce tailored recommendations; in a user-product bipartite rating network, using a few manually identified fraudulent user accounts, SSL is useful to identify other fraudulent accounts [4, 10]; SSL can identify protein functions from networks of their physical interaction using just a few labels [31].

Traditional graph SSL algorithms leverage a key property of real-world networks: the homophily of vertices [5, 20], i.e., the nearby vertices in a graph are likely to have the same label. However, these methods tend to be limited by the fact that all the neighbors of a vertex are not equal. Consider your own friendship network where you have many acquaintances, but only a few close friends. In fact, prior research has shown that vertices with a strong connection participate in several higher-order structures, such as dense subgraphs and cliques [12–14, 26]. Thus, leveraging the higher-order structure between vertices is crucial to accurately label the vertices.

Let us elaborate this using a small friendship network example, shown in Figure 1. The central vertex, Alice, participates in a closely-knit community with three friends B, C, and D, all of whom know each other. In addition, she has four acquaintances P, Q, R, and S from different walks of her life. Let the vertices be labeled by their political beliefs—vertices B, C, and D have the same blue label; and the rest of the vertices have the red label. Even though Alice has more red connections than blue, the connection between Alice, B, C, and D is stronger as Alice participates in three 3-cliques and one 4-clique with them. In contrast, Alice has no 3- and 4-cliques with P, Q, R, and S. Owing to the stronger connection with the red nodes, Alice should be labeled red as well. However, traditional graph SSL techniques that rely on edges alone label Alice as blue [38, 39]. This calls for graph SSL methods that look beyond edges to leverage the signal present in higher-order structures to label vertices.

Figure 1: Graph SSL approaches which take only edges into account incorrectly classify the unlabeled central vertex ‘Alice’ as blue. By leveraging higher-order network structures, the proposed HOLS correctly labels Alice as red.

Our present work focuses on three key research questions:

• **RQ1.** How do the data reveal that higher-order network structures are homogeneous in labels?
• **RQ2.** How can we leverage higher-order network structures for graph SSL in a principled manner?
• **RQ3.** Do higher-order structures help improve graph SSL?

Accordingly, our contributions can be summarized as follows: (i) Analysis: We conduct an empirical analysis on four diverse real-world networks to confirm that label homogeneity is indeed a highly prevalent characteristic in higher-order structures. (ii) Algorithm: We develop Higher-Order Label Spreading (HOLS) to leverage higher-order structures during graph semi-supervised learning. HOLS works for any user-inputted higher-order structure and in the base case, is equivalent to the standard edge-based label spreading [38]. (iii) Effectiveness: Via experiments, we show that label spreading via higher-order structures strictly outperforms label
spreading via edges. Notably, HOLS outperforms state-of-the-art deep learning methods by a statistically significant margin across most datasets, while running 15× less time faster on a stock laptop.

For reproducibility, all the code and datasets are available at https://github.com/dhivyaeswaran/hols.

2 RELATED WORK

Traditional graph SSL approaches: By far, the most widely used graph SSL techniques are label propagation [39] and label spreading [38]. Label propagation (LP) clamps labeled vertices to their provided values and uses a graph Laplacian regularization, while label spreading (LS) uses a squared Euclidean penalty as supervised loss and normalized graph Laplacian regularization which is known to be better-behaved and more robust to noise [32]. Both these techniques permit closed-form solution and are extremely fast in practice, scaling well to billion-scale graphs. Consequently, a number of techniques build on top of these approaches, for example, to allow inductive generalization [6, 33], to incorporate certainty [29], and so on. When the graph is viewed as pairwise Markov random field, belief propagation (BP) [36] may be used to recover the exact marginals on the vertices. BP can handle network effects beyond just homophily; however, it has well-known convergence problems from a practitioner’s point of view [25]. While traditional techniques, in general, show many desirable theoretical properties such as closed-form solution, convergence guarantees, connections to spectral graph theory [39] and statistical physics [36], as such, they do not account for higher-order network structures.

Recent graph SSL approaches differ from traditional SSL methods in training embeddings of vertices to jointly predict labels as well as the neighborhood context in the graph. Specifically, Planetoid [35] uses skipgrams, while GCN [17] uses approximate spectral convolutions to incorporate neighborhood information. MixHop [1] can learn a general class of neighborhood mixing functions for graph SSL. As such, these do not incorporate specific higher-order structures provided by the user. Further, their performance in practice tends to be limited by the availability of ‘good’ vertex features for initializing the optimization procedure.

Hybrid approaches for graph SSL: Another way to tackle the graph SSL problem is a hybrid approach to first extract vertex embeddings using an unsupervised approach such as node2vec [11], DeepWalk [21] or LINE [30] and then use the available labels to learn a transductive classifier such as an SVM [16]. Such methods, however, neither have well-understood theoretical properties nor do they optimize for a single objective in an end-to-end manner.

Comparison: We compare the best performing HOLS algorithm (HOLS-3 which uses triangles in addition to edges) qualitatively to prominent SSL approaches in Table 1 and quantitatively via experiments to representative methods from the above categories: LP and LS (traditional), GCN (recent) and node2vec + TSVM (hybrid).

3 MOTIVATION

Recent work has shown that graphs from diverse domains have many striking higher-order network structures [8] which can be leveraged to improve tasks such as clustering [37], link prediction [2, 7] and ranking [23]. From a graph SSL point of view, the explicit use of higher-order network structures has remained limited to belief propagation over 2 × 2 image cliques to improve image denoising, segmentation and rendering in computer vision [18, 22]. However, till date, the importance of higher-order network structures in SSL atop explicit graph data has largely remained unexplored.

In this section, we motivate the need to consider higher-order network structures for semi-supervised learning through empirical analysis of four diverse large real-world networks. As we will see, higher-order label homogeneity—i.e., the tendency of vertices participating in higher-order structures to share the same label—is remarkably more common than expected in real-world graphs.

3.1 Dataset Description

We use following network datasets for our empirical analysis and experiments. Table 2 summarizes some important dataset statistics.

- **EUEmail** [19] is an e-mail communication network from a large European research institution. Vertices indicate members of the institution and an edge between a pair of members indicates that they exchanged at least one email. Vertex labels indicate membership to one of the 42 departments.
- **PolBlogs** [3] is a network of hyperlinks between blogs about US politics during the period preceding the 2004 presidential election. Blogs are labeled as right-leaning or left-leaning.
- **Cora** [27] is a citation network among papers published at computer science conferences. Vertex labels indicate membership of the paper to one of 10 areas that the paper belongs to based on its venue of publication, e.g., Artificial Intelligence, Databases, Networking.
- **Pokec** [28] is the most popular online social network in Slovakia. Vertices indicate users and edges indicate friendships. From the furnished user profile information, we extract the locality or ‘kraj’ that users belong to and use them as labels.

These datasets exhibit homophily [5, 20]: people typically e-mail others within the same department; blogs tend to link to others
having the same political leaning, papers mostly cite those from the same area; people belonging to the same locality are more likely to meet and become friends. In all cases, we omit self-loops and take the edges as undirected and unweighted.

3.2 Higher-Order Label Homogeneity

We now examine the label homogeneity of higher-order structures in real-world networks. We focus on cliques, which are quintessential dense subgraphs which form the essential building blocks of many networks [12–14, 26], moreover, recent advancements enable fast and efficient counting and enumeration of cliques [9, 15]. We will stick to $k \in \{2, 3, 4, 5\}$ for computational reasons.

**Label configuration.** We introduce the term label configuration to capture a function of vertex-label assignments that is invariant under the permutation of vertices and labels. For example, a 2-clique has two label configurations: ‘2’ where both incident vertices have the same label and ‘1-1’ where they have different labels. A 3-clique has three label configurations: ‘3’ where all vertices have the same label, ‘2-1’ where two of them share the same label and third vertex has a different label and ‘1-1-1’ where each vertex has a different label. Similarly, a 4-clique has 5 label configurations (4, 3-1, 2-2, 2-1-1, 1-1-1-1) and a 5-clique has 7 label configurations (5, 4-1, 3-2, 3-1-1, 2-2-1, 2-1-1-1, 1-1-1-1-1). Note that not all label configurations may be possible (e.g., 1-1-1 is impossible for a triangle in a 2-class problem) and still fewer may actually occur in practice. Thus, as a first step toward analyzing label homogeneity, we study the distribution over $k$-clique label configurations.

Figure 2 plots the observed distribution over $k$-clique label configurations (in orange). For ease of interpretation, we compare this to the expected distribution of a random baseline (in green) which shuffles vertex labels fixing the graph structure. In each subfigure, more homogeneous label configurations where many vertices share the same label are displayed on the left, whereas less homogeneous label configurations where vertices have different labels are shown towards the right. We observe that, in all cases, the most homogeneous label configuration is $3 \sim 1000 \times$ more common than expected, while the least homogeneous (and possible) label configuration is $5 \sim 100 \times$ rarer than expected. While quantifying label homogeneity in higher-order structures is left to future work, Figure 2 establishes concrete evidence that real-world $k$-cliques are indeed more homogeneous in labels than can be explained by random chance.

We now show how this observation paves the way for more accurate semi-supervised learning using higher-order structures.

4 ALGORITHM

In the section, we derive the proposed higher-order label spreading algorithm and shows its desirable theoretical properties.
4.1 Notation
Let $G = (V, E)$ be a graph with vertex set $V$ and edge set $E$. Edges are undirected with $w_{ij}$ representing the edge weight between vertices $i, j \in V$. Each vertex has a unique label $\ell(i) \in \{1, 2, \ldots, C\}$.

Let $\mathcal{K}$ be the set of network structures or motifs (e.g., edges, triangles) which we want to leverage for graph semi-supervised learning. For a given motif $\kappa \in \mathcal{K}$, let $|\kappa|$ denote its size which is its number of participating vertices. For example, when $\kappa$ is a triangle, $|\kappa| = 3$. Further, suppose $Q_k$ is the set of occurrences of a motif $\kappa$ in graph $G$ and each such occurrence $q \in Q_k$ has a weight $w_q$ (e.g., computed as the product of incident edge weights). Use $\mathbb{I} \{ \cdot \}$ to denote the indicator function, which evaluates to 1 when the enclosed expression is true. For example, $\mathbb{I} \{ i \in q \}$ is one if the vertex $i$ is part of the subgraph $q$.

4.2 Generalized Loss Function
Let $y_i \in \{0, 1\}^C$ be the one-hot vector of provided label for a labeled vertex $i$ such that $y_{ic} = 1$ if vertex $i$ has a label $c$ and is zero otherwise. We propose to infer the final labels $x_i \in \{0, 1\}^C$ (where $\sum_c x_{ic} = 1 \forall i \in V$) by minimizing the following loss function:

$$L = (1 - \eta)L_s + \eta L_g = (1 - \eta)\sum_{i \in V} \alpha_{\kappa}L_{g,\kappa}$$

where $L_s = \frac{1}{2} \sum_{i,j} ||x_i - x_j||^2$ is the supervised loss, which penalizes the deviation of inferred labels from their provided values, while $L_{g,\kappa}$ is the graph loss with respect to motif $\kappa$, which enforces the inferred labels to be ‘smooth’ over all occurrences of $\kappa$ as:

$$L_{g,\kappa} = \frac{1}{2} \sum_{q \in Q_k} \sum_{i,j \in q} w_q \mathbb{I} \{ i \in q \land j \in q \} ||x_i - x_j||^2$$

A parameter $\eta \in (0, 1)$ trades off supervised and graph losses, while $\alpha_{\kappa} \in (0, 1)$ captures the importance weight of $\kappa$ in semi-supervised learning. Note $\sum_{\kappa \in \mathcal{K}} \alpha_{\kappa} = 1$. Now, define $\kappa$-participation matrix as $W^{(\kappa)} = [w^{(\kappa)}_{ij}]$ where each entry $w^{(\kappa)}_{ij}$ denotes the total weight of $\kappa$-motifs that vertices $i$ and $j$ participate in. We have:

$$w^{(\kappa)}_{ij} = \sum_{q \in Q_k} w_q \mathbb{I} \{ i \in q \land j \in q \}$$

Observe that each pairwise loss term $||x_i - x_j||^2$ in Equation (1) appears with a total weight $w^{(\kappa)}_{ij}$ given by $w^{(\kappa)}_{ij} = \sum_{\kappa \in \mathcal{K}} \alpha_{\kappa}w^{(\kappa)}_{ij}$ using which we may simplify the graph loss as:

$$L_g = \frac{\eta}{2} \sum_{i,j} w^{(\kappa)}_{ij} ||x_i - x_j||^2$$

Thus, Equation (4) establishes that the graph loss from Equation (1) is equivalent to that of label propagation [39] on a modified graph with adjacency matrix $W' = \sum_{\kappa \in \mathcal{K}} \alpha_{\kappa}W^{(\kappa)}$ where each edge of the original graph has been re-weighted according to the total weight of $\kappa$-motifs it participates in, scaled by its importance $\alpha_{\kappa}$, and finally summed over all motifs $\kappa \in \mathcal{K}$ of interest. We will use this connection to derive a closed-form solution to HOLS.

4.3 Closed-Form and Iterative Solutions
Let $Y = [y_1, \ldots, y_N]^T$ and $X = [x_1, \ldots, x_N]^T$ be the $N \times C$ matrices of prior and inferred labels where $|V| = N$ is the total number of vertices. Let $D' = [d'_{ij}]$ be the diagonal degree matrix for the modified graph adjacency $W' = [w'_{ij}]$. Thus, $d'_{ii} = \sum_{j \neq i} w'_{ij}$ and $d'_{ij} = 0$ if $i \neq j$. Let $L' = D' - W'$ be the Laplacian matrix for the modified graph. Equation (4) can be re-written in matrix format as:

$$L = \frac{1 - \eta}{2} ||X - Y||^2_F + \frac{\eta}{2} X^T L' X$$

We also consider a version of the loss function which uses the normalized Laplacian $L' = D'^{-1/2}L'D'^{-1/2}$ for regularization:

$$L = \frac{1 - \eta}{2} ||X - Y||^2_F + \frac{\eta}{2} X^T L' X$$

Using $L'$ in place of $L$ performs as well if not better in practice; and moreover provides certain theoretical guarantees (see Proposition 4.2, and also [32]). Therefore, we will use Equation (6) as the loss function for our higher-order label spreading and refer to it as $L_{HOLS}$. A closed-form solution can now be obtained by differentiating $L_{HOLS}$ with respect to $X$ and setting it to zero. Thus:

$$X = (1 - \eta)(1 - \eta(I - L'))^{-1} Y$$

Thus, using Equation (7), we are able to compute the optimal solution to HOLS, as long as the inverse of $I - \eta(I - L')$ exists. Due to the use of normalized Laplacian regularization, the following holds:

**Proposition 4.1 (Generalized Label Spreading).** The proposed HOLS algorithm reduces to traditional label spreading [38] for the base case of using only edge motifs, i.e., $\mathcal{K} = \{K_2\}$.

This generalization grants HOLS its name. In practice, matrix inversion is computationally intensive and tends to be numerically unstable. Hence, we resort to an iterative approach to solve Equation (7) by first initializing $X$ to an arbitrary value and then repeatedly applying the following update:

$$X \leftarrow \eta(I - L')X + (1 - \eta)Y$$
Proposition 4.2 describes the theoretical properties of this approach.

**Proposition 4.2 (Convergence Guarantee for HOLS).** The iterative update in Equation (8) always converges to the unique fixed point given in Equation (7) for any choice of initial $X$.

This can be proved using the theory of sparse linear systems [24]. The overall algorithm of HOLS is summarized in Algorithm (1). First, for each motif $k \in K$, construct its $k$-participation matrix by enumerating all its occurrences. Note that the enumerated occurrences are processed one by one on the fly to update the participation matrix and discarded (no need for storage). Moreover, the enumeration for different motifs can be done in parallel. The participation matrices are combined into a single modified graph adjacency $W'$; applying the iterative updates from Equation (8) finally results in labels for the unlabeled vertices. In practice, the iterative updates are applied until entries in $X$ do not change up to a precision $\epsilon$ or until a maximum number of iterations $T$ is reached.

### 4.4 Complexity Analysis

When only cliques are used as motifs $K$ for semi-supervised learning, the following space and time complexity bounds hold:

**Proposition 4.3 (Space Complexity of HOLS).** The space complexity of HOLS for a graph with $N$ vertices, $M$ edges and $C$ classes is $O(M + NC)$ independent of motif size and number of motifs used, provided all motifs are cliques.

The proof follows by observing that the modified adjacency $W'$ is at least as sparse as the original adjacency matrix $W$ when the motifs used are cliques. Moreover, $W'$ can be constructed on the fly by enumerating—but not storing—the occurrences of all motifs.

**Proposition 4.4 (Time Complexity of HOLS).** The time complexity of HOLS over a graph with $M$ edges, $C$ classes and a degeneracy (core number) of $k_{\text{max}}$ using $K = \{K_2, \ldots, K_n\}$ is given by

$$O\left(M \sum_{k=2}^{k_{\text{max}}} k \left( \frac{k_{\text{max}}}{2} \right)^{k-3}\right)$$

for the construction of $K_k$-participation matrices plus $O(MC)$ per iterative update using Equation (8).

The proof follows from Theorem 5.7 of [9].

Despite the exponential complexity in $k$, we are able to enumerate cliques quickly using the sequential kClist algorithm [9]. For example, our largest Pokec dataset has 21M edges, 32M triangles, 43M 4-cliques and 53M 5-cliques; and the enumeration of each took a maximum of 20 seconds on a stock laptop. Thus, HOLS remains fast and scalable when reasonably small cliques are used. Further, as we show in experiments, using triangles (3-cliques) in addition to edges typically suffices to achieve the best classification performance across a wide range of datasets.

### 5 EXPERIMENTS

We empirically evaluate the proposed algorithm on the four real-world network datasets described in Section 3.1.

#### 5.1 Experimental Setup

We implement HOLS in MATLAB and run the experiments on MacOS with 2.7 GHz Intel Core i5 processor and 16 GB main memory.

**Baselines.** We compare HOLS to the following baselines: (1) Label Propagation (LP) [39] which uses Laplacian regularization. (2) Label Spreading (LS) [38] which uses normalized Laplacian regularization. (3) node2vec+TSVM which generates unsupervised vertex embeddings using node2vec [11] and learns decision boundaries in the embedding space using a transductive SVMs [16]. (4) Graph Convolutional Network (GCN) [17] which is an end-to-end semi-supervised learner using neural networks. We implement LP and LS in MATLAB, and use open-sourced code for the rest.

**Parameters.** By default, we use a weight of $\eta = 0.5$ for supervised loss and $K = \{K_2, K_3\}$ motifs (edges and triangles) for HOLS. The importance weight for triangles $\alpha_{k_3}$ is tuned in $\{0.1, 0.2, \ldots, 0.9\}$ for each dataset and results are reported on the best performing value. We use $\eta = 0.5$ for LS as well. LP, LS and HOLS are run until labels converge to a precision of $\epsilon$ or until $T$ iterations are completed, whichever occurs sooner. We set $\epsilon = 10^{-6}$ and $T = 500$. We use the default hyperparameters for GCN, node2vec and TSVM. We supply 100, 20, 100 and 1000 labels for EuEmail, PolBlogs, Cora and Pokec datasets, where the vertices to label are chosen by stratified sampling based on class. These correspond to label fractions of 5%, 1.6%, 0.4% and 0.06% and on an average, 1, 10, 10 and 100 labeled vertices per class respectively.

**Evaluation metrics.** We quantify success using accuracy, which is the fraction of unlabeled vertices which are correctly classified. We also note down the end-to-end running time for all computation including any pre-processing such as clique enumeration, but excluding I/O operations. All reported values are averaged over five runs, each run differing in the set of labeled vertices.

#### 5.2 Results

**Accuracy:** Accuracies of all methods are summarized in Table 3 (left). The values for node2vec+TSVM on Pokec dataset are missing as the method did not terminate within 24 hours (’T.L.E.’).

First, we observe in Table 3 (left) that HOLS consistently leads to statistically significant improvements over LS, showing that using higher-order structures for label spreading helps. In addition, HOLS outperforms all baselines in three out of four datasets. The improvements over the best baseline are statistically significant ($p = 0.05$) according to a two-sided micro-sign test [34] in at least three out of five runs. Interestingly, for the smaller datasets (EuEMAIL and PolBlogs), while GCN outperforms LS, GCN loses to HOLS when triangles are used. node2vec+TSVM performs slightly better than HOLS on PolBlogs, however, the increase over HOLS is not statistically significant. For the larger datasets with $< 0.5\%$ labeled vertices, HOLS performs the best and LS follows closely.

**Running Time:** The running time of HOLS and all the baselines is summarized in Table 3 (right). Notably, we see that HOLS runs in less than 2 minutes for graphs with over 21 million edges (the Pokec graph), demonstrating its real-world practical scalability.

We observe that LS is the fastest of all methods and HOLS comes a close second for three out of four datasets. The small difference in running time predominantly stems from the construction of triangle participation matrix. Furthermore, HOLS is over 15× faster than the recent GCN and node2vec+TSVM baselines, for comparable and often better values of accuracy.

**Accuracy vs. Maximum Clique Size:** Fixing the motif set as $K = \{K_2, K_3, \ldots, K_k\}$, we vary $k = 2, 3, 4, 5$ to study the marginal benefit of including higher-order cliques in graph SSL. The motif
Table 3: Accuracy and Running Time (averaged over five runs): In each column, the best value is bold and underlined, and the second best is underlined. Asterisk (*) denotes statistically significant differences (p < 0.05) compared to the second best.

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<th>Method</th>
<th>Metric</th>
<th>EuEmail</th>
<th>PolBlogs</th>
<th>CORA</th>
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Figure 3: Variation of accuracy with (a) the maximum clique size $k$ used, and (b) importance weight $\alpha_{K_3}$ to triangles.

Figure 4: Case studies from PolBlogs dataset showing extended ego-networks of vertices v702 (left) and v1153 (right) which are both incorrectly classified by LS but correctly classified when triangles are taken into account using HOLS.

weights are tuned in $\alpha_{K_3} \in \{0, 0.1, \ldots, 0.9\}$, ensuring that edges are given a weight $\alpha_{K_3} \geq 0.1$ for a connected graph, and further, all motif weights sum to 1. The best performing motif weights were used to generate Figure 3(a), which plots the relative improvement in accuracy over LS that uses edges only. We note that label spreading via higher-order structures strictly outperforms label spreading via edges. The gain is the most when using 3-cliques (triangles). Subsequent higher-order cliques did not lead to additional performance gain in most cases, presumably because their occurrences tend to be concentrated around a few high-degree vertices.

Accuracy vs. Importance Weight To Triangles: Fixing the motif set to $\mathcal{K} = \{K_2, K_3\}$, we vary the importance weight $\alpha_{K_3}$ to triangles in $\{0, 0.1, \ldots, 0.9\}$ to understand its effect on accuracy. Figure 3(b) shows that the accuracy gain of HOLS over LS increases with an increase in triangle weight for most graphs. The only exception is CORA, where the accuracy gain grows until $\alpha_{K_3} = 0.4$ before decreasing and eventually turning negative. Overall, triangles consistently help over a large range of motif weights.

Case Studies: In Figure 4, we look at real examples from the PolBlogs dataset to dig deeper into when HOLS improves over LS. Question mark denotes the central vertices v702 and v1153 of interest with ground truth labels ‘blue’ and ‘red’ respectively. The direct neighbors of the both vertices are unlabeled and a few second hop neighbors are labeled with one of two labels: ‘blue’ or ‘red’.

In both cases, LS correctly labels all vertices except the central vertex. This is because LS utilizes only edge-level information: v702 has more second-hop red neighbors than blue whereas v1153 has more second-hop blue neighbors than red. On the other hand, HOLS leverages the fact that v702 participates in one triangle with (inferred) blue vertices and that v1153 participates in several triangles with (inferred) red vertices and thus produces the correct labels.

6 CONCLUSION

In this paper, we demonstrated that label homogeneity—the tendency of vertices participating in a higher-order structure to share the same label—is a prevalent characteristic in real-world graphs. We created an algorithm to exploit the signal present in higher-order structures for more accurate semi-supervised learning over graphs. Experiments on real-world data showed that using triangles along with edges for label spreading leads to statistically significant accuracy gains compared to the use of edges alone.

This work opens the avenue for several exciting future research directions. First, we need principled measures quantifying label homogeneity to aid comparison across diverse graphs and higher-order structures. Next, having seen the improvements in real-world graphs, it becomes fundamental to understand the benefits in random graph models. Finally, it is crucial to develop algorithms exploiting higher-order structures which can cater to the increasingly heterogeneous and dynamic nature of real-world graphs at scale.

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